## ECE-250 – Algorithms and Data Structures (Winter 2012) Tutorial 2 (2012-01-26)

 ${\bf 1}$  - Assuming well-behaved positive functions such as those that we encounter when analyzing algorithms, show that:

- (a) Little-oh defines a partial ordering.
- (b) Big-Oh defines a weak ordering.
- (c) Big-Theta is an equivalence relation (one of the three conditions was already done during last week's tutorial session).
- **2** Determine the run time (worst-case only) of the following operations:

(a) Find the first negative value in an array of:

- (i) Random values
- (ii) Ordered values in ascending sequence
- (iii) Ordered values in descending sequence
- (b) Same as part (a), but for a linked list

(c) Insert an element into a singly-linked list with random values:

- (i) After a given element (i.e., you're given a pointer to a node)
- (ii) Before a given element
- (iii) Exactly at the middle (plus/minus rounding if the list contains an odd number of elements)

**3** - Compare (a discussion suffices; no need to formally determine the actual figures) the run times (if applicable, discuss the "actual time" performance) of operations on a queue (enqueue and dequeue) when we implement it in terms of:

- (a) An array (linear array)
- (b) A circular array or circular buffer
- (c) A linked list

**4** - (If enough time) Determine the run time of the following function, operating on a polygon (a sequence of points) represented as an array of point objects (we have some class Point to represent the points/vertices of the polygon):

The function needs to check a condition for every pair of points; the condition is symmetric (i.e., if the condition is met for points  $p_i, p_j$ , then it is also met for the pair  $p_j, p_i$ ), and it is checked in constant time. It is also known that the condition is never met between a point and the five points that follow.

Determine the runtime of a function meeting the above requirements.