

ECE 602 – Introduction to Optimization

Solutions to Home Assignment 1

Exercise 1

First, we rewrite the cost function as

$$f(x) = \langle \mathbf{1}, \sqrt{(Ax)^2 + \epsilon} \rangle,$$

with $\mathbf{1} = [1, 1, \dots, 1]^T \in \mathbf{R}^n$, where the functions $(\cdot)^2$ and $\sqrt{\cdot}$ are applied element-wise. In this case,

$$\begin{aligned} df(x) &= f(x + dx) - f(x) = \langle \mathbf{1}, \sqrt{(A(x + dx))^2 + \epsilon} \rangle - f(x) = \\ &= \langle \mathbf{1}, \sqrt{(Ax + Adx)^2 + \epsilon} \rangle - f(x) = \langle \mathbf{1}, \sqrt{(Ax)^2 + 2(Ax) \cdot (Adx) + (Adx)^2 + \epsilon} \rangle - f(x) \approx \\ &\approx \langle \mathbf{1}, \sqrt{((Ax)^2 + \epsilon) + 2(Ax) \cdot (Adx)} \rangle - f(x), \end{aligned}$$

where we discard the high-order term $(Ax)^2$ due to its relative smallness. Note that, in the expressions above, the dot \cdot denotes an element-wise product of two vectors.

Next, we use the following (1st-order) Taylor approximation around $((Ax)^2 + \epsilon)$

$$\begin{aligned} \sqrt{((Ax)^2 + \epsilon) + 2(Ax) \cdot (Adx)} &\approx \sqrt{(Ax)^2 + \epsilon} + \frac{1}{\sqrt{(Ax)^2 + \epsilon}} \cdot (Ax) \cdot (Adx) = \\ &= \sqrt{(Ax)^2 + \epsilon} + \frac{Ax}{\sqrt{(Ax)^2 + \epsilon}} \cdot (Adx), \end{aligned}$$

where the vector division is assumed to be element-wise. Thus, we have

$$\begin{aligned} df(x) &= \langle \mathbf{1}, \sqrt{(Ax)^2 + \epsilon} + \frac{Ax}{\sqrt{(Ax)^2 + \epsilon}} \cdot (Adx) \rangle - f(x) = \\ &= \langle \mathbf{1}, \sqrt{(Ax)^2 + \epsilon} \rangle + \langle \mathbf{1}, \frac{Ax}{\sqrt{(Ax)^2 + \epsilon}} \cdot (Adx) \rangle - f(x) = \langle \mathbf{1}, \frac{Ax}{\sqrt{(Ax)^2 + \epsilon}} \cdot (Adx) \rangle = \end{aligned}$$

$$= \left\langle \frac{Ax}{\sqrt{(Ax)^2 + \epsilon}}, Adx \right\rangle = \left\langle A^T \left(\frac{Ax}{\sqrt{(Ax)^2 + \epsilon}} \right), dx \right\rangle.$$

Therefore,

$$\nabla f(x) = A^T \left(\frac{Ax}{\sqrt{(Ax)^2 + \epsilon}} \right).$$

Exercise 2

PART A

- a) Let $x_1, x_2 \in C_\alpha$. Let's show that, for any $0 \leq \theta \leq 1$, $\theta x_1 + (1 - \theta)x_2 \in C_\alpha$, meaning $f(\theta x_1 + (1 - \theta)x_2) \leq \alpha$. Indeed, due to the convexity of f , we have

$$f(\theta x_1 + (1 - \theta)x_2) \leq \theta f(x_1) + (1 - \theta)f(x_2) \leq \theta \alpha + (1 - \theta)\alpha = \alpha.$$

- b) Let $Q_1, Q_2 \in \mathbf{S}_+^n$. Let's show that, for any $0 \leq \theta \leq 1$, $Q = \theta Q_1 + (1 - \theta)Q_2 \in \mathbf{S}_+^n$, meaning $x^T Q x \geq 0$, for any $x \in \mathbf{R}^n$. Indeed, we have

$$x^T Q x = x^T (\theta Q_1 + (1 - \theta)Q_2) x = \theta x^T Q_1 x + (1 - \theta) x^T Q_2 x \geq 0,$$

since both $x^T Q_1 x \geq 0$ and $x^T Q_2 x \geq 0$.

PART B

- a) The Hessian of $f(x) = 1/2 x^T Q x + c^T x$ is equal to Q , which is a positive-definite matrix. Therefore, the function is convex (albeit, not strictly).
- b) Let $x_1, x_2 \in \mathbf{R}$. To prove the convexity of $f(x) = g(h(x))$, we need to show that, for any $\theta \in [0, 1]$, we have

$$f(\theta x_1 + (1 - \theta) x_2) \leq \theta f(x_1) + (1 - \theta)f(x_2).$$

First, we observe that, by the convexity of h , we have

$$h(\theta x_1 + (1 - \theta) x_2) \leq \theta h(x_1) + (1 - \theta)h(x_2).$$

Consequently, due to the monotonicity of g , we have

$$g(h(\theta x_1 + (1 - \theta) x_2)) \leq g(\theta h(x_1) + (1 - \theta)h(x_2)).$$

At the same time, due to the convexity of g , we can write

$$g(\theta h(x_1) + (1 - \theta)h(x_2)) \leq \theta g(h(x_1)) + (1 - \theta)g(h(x_2)).$$

In summary,

$$f(x) = g(\theta h(x_1) + (1 - \theta)h(x_2)) \leq \theta g(h(x_1)) + (1 - \theta)g(h(x_2)) = \theta f(x_1) + (1 - \theta)f(x_2).$$

Exercise 3

Let p^* be the global minimum of f . Then the set of all global minimizers of f can be defined as

$$C_{p^*} = \{x \in \mathbf{R}^n \mid f(x) \leq p^*\}.$$

This is a sublevel set of a convex function, which makes it convex.

Exercise 4

a) To show that $\|z\|_\infty = \sup_{\|x\|_1 \leq 1} x^T z$, we first notice that

$$x^T z \leq |x^T z| \leq \sum_{i=1}^n |x_i| |z_i| \leq \max_i |z_i| \sum_{i=1}^n |x_i| = \|z\|_\infty \|x\|_1 \Big|_{\|x\|_1=1} = \|z\|_\infty.$$

However, it is not clear yet if the bound is tight. To show that, we need to find a vector x (with $\|x\|_1 = 1$) for which $x^T z = \|z\|_\infty$. Such x can be defined as

$$x_i = \begin{cases} \text{sign}(z_{i^*}), & \text{if } i = i^* \\ 0, & \text{otherwise} \end{cases},$$

with $i^* = \arg \max_{1 \leq i \leq n} |z_i|$.

b) To show that $\|z\|_2 = \sup_{\|x\|_2 \leq 1} x^T z$, we simply use the Cauchy-Schwarz inequality

$$x^T z \leq |x^T z| \leq \|z\|_2 \|x\|_2 \Big|_{\|x\|_2=1} = \|z\|_2,$$

which is known to be tight for $x = z/\|z\|_2$.