

# Optimal Power Allocation And Network Beamforming For OFDM-Based Relay Networks

ECE 602 Final Project- Winter 2019

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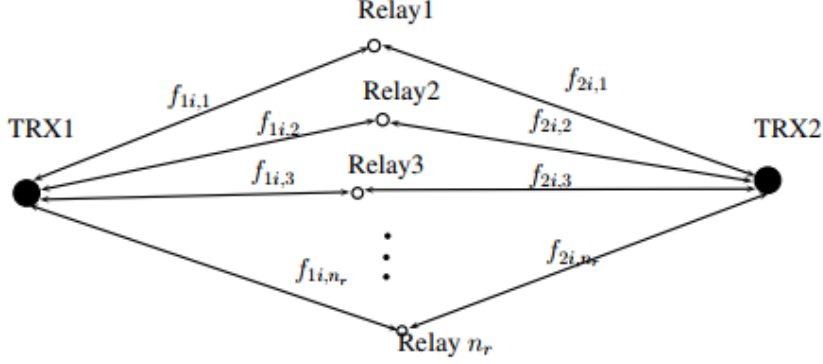
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## 1. Problem formulation

Authors of the paper have investigated the problem of optimal power allocation in an OFDM two-way relay network with an arbitrary number of relays. They have assumed that two way relaying is performed employing analog network coding, accordingly, the optimal power allocation across subcarriers and among a relay and two communication nodes were obtained by minimizing the total power consumption in the network subject to two separate rate constraints for each transceiver. An efficient algorithm was proposed by the authors of the proposed optimization problem. Based on the paper, a two-way relay network consisting of a pair of transceivers (TRX1 and TRX2) and  $n_r$  relay nodes, as shown in Fig.1[1]. All nodes use OFDM for data transmission over  $n_c$  subcarriers. Authors assume the same set of subcarriers is used for the relays to receive and re-transmit received signal from TRX1 and TRX2, and thus, all channels to and from each relay are reciprocal. Let  $p_{1i}$  and  $p_{2i}$  denote the transmit powers allocated to the  $i$ th subcarrier at TRX1 and TRX2, respectively, where  $i \in \{1, 2, \dots, n_c\}$ . We can write the  $(n_r \times 1)$  complex vector  $x_i$  of the signals received by the relays over  $i$ th subcarrier as:

$$x_i = \sqrt{p_{1i}} f_{1i} s_{1i} + \sqrt{p_{2i}} f_{2i} s_{2i} + \nu_i \quad (1)$$

where  $\nu_i$  is an  $(n_r \times 1)$  complex vector representing the relay noises on the  $i$ th subcarrier,  $s_{1i}$  ( $s_{2i}$ ) is the information symbol transmitted by TRX1(TRX2) over the  $i$ th subcarrier, and  $f_{1i}$  ( $f_{2i}$ ) is the  $(n_r \times 1)$  complex vector of the channel coefficients between TRX1(TRX2) and the relays corresponding to the  $i$ th subcarrier.[1]



**Fig. 1.** A two-way relay network.

## 2. Proposed solution

The authors of the paper aim to minimize the total transmit power subject to two separate constraints on the rates of the end users. The proposed optimization problem is presented as below:

$$\begin{aligned}
 & \min_{\mathbf{p}_1, \mathbf{p}_2, \mathcal{W}} && P_T(\mathbf{p}_1, \mathbf{p}_2, \mathcal{W}) \\
 & \text{subject to} && \frac{1}{2} \sum_{i=1}^{n_c} \log(1 + \text{SNR}_{1i}(p_{2i}, w_i)) \geq r_1^{\max} \\
 & && \frac{1}{2} \sum_{i=1}^{n_c} \log(1 + \text{SNR}_{2i}(p_{1i}, w_i)) \geq r_2^{\max}
 \end{aligned} \tag{2}$$

where  $p_1 = [p_{11} \ p_{12} \ \dots \ p_{1n_c}]^T$ ,  $p_2 = [p_{21} \ p_{22} \ \dots \ p_{2n_c}]^T$ ,  $\mathcal{W} = \{w_i\}_{i=1}^{n_c}$ .  $r_1^{\max}$  and  $r_2^{\max}$  are the minimum required sum-rate for TRX1 and TRX2, respectively. Also, the  $i$ th subcarrier SNRs at TRX1 and TRX2 and the total consumed power in the network are respectively given by equation (3) and(4):

$$\text{SNR}_{1i}(p_{2i}, w_i) = \frac{p_{2i} w_i^H h_i h_i^H w_i}{1 + w_i^H D_1 w_i}$$

$$\text{SNR}_{2i}(p_{1i}, w_i) = \frac{p_{1i} w_i^H h_i h_i^H w_i}{1 + w_i^H D_2 w_i} \quad (3)$$

$$P_T(p_1, p_2, W) = \sum_{i=1}^{n_c} p_{1i} + p_{2i} + p_{r,i} \quad (4)$$

where  $p_{r,i}$  is the total relay power consumed over the  $i$ th subcarrier, can be written as

$$p_{r,i} = w_i^H (p_{1i} D_1 + p_{2i} D_1 + I) w_i \quad (5)$$

It can be easily shown that, at the optimum, the two rate constraints in (2) must be satisfied with equality. To solve (2), we can rewrite it as:

$$\begin{aligned} \min_{\mathbf{r}_1, \mathbf{r}_2} \quad & \sum_{i=1}^{n_c} \min_{p_{1i}, p_{2i}, \mathbf{w}_i} p_{1i} + p_{2i} + p_{r,i}(p_{1i}, p_{2i}, \mathbf{w}_i) \\ \text{subject to} \quad & \frac{1}{2} \log(1 + \text{SNR}_{1i}(p_{2i}, \mathbf{w}_i)) \geq r_{1i} \\ & \frac{1}{2} \log(1 + \text{SNR}_{2i}(p_{1i}, \mathbf{w}_i)) \geq r_{2i} \\ & \mathbf{1}^T \mathbf{r}_1 = r_1^{\max}, \mathbf{r}_1 \succcurlyeq \mathbf{0} \\ & \mathbf{1}^T \mathbf{r}_2 = r_2^{\max}, \mathbf{r}_2 \succcurlyeq \mathbf{0} \end{aligned} \quad (6)$$

where the auxiliary variable  $r_{1i}$  and  $r_{2i}$  are rates of TRX1 and TRX2, respectively, over the  $i$ th subcarrier and we define  $r_j = [r_{j1} \ r_{j2} \ \dots \ r_{jn_c}]^T$ , for  $j = 1, 2$ . the above optimization problem(6) is converted to the following from based on the paper[1]:

$$\begin{aligned} \min_{\mathbf{r}_1, \mathbf{r}_2, \mathcal{W}} \quad & \sum_{i=1}^{n_c} (2^{2r_{1i}} + 2^{2r_{2i}} - 2) \xi_i(\mathbf{w}_i) + \mathbf{w}_i^H \mathbf{w}_i \\ \text{subject to} \quad & \mathbf{1}^T \mathbf{r}_1 = r_1^{\max}, \mathbf{r}_1 \succcurlyeq \mathbf{0} \\ & \mathbf{1}^T \mathbf{r}_2 = r_2^{\max}, \mathbf{r}_2 \succcurlyeq \mathbf{0}. \end{aligned} \quad (7)$$

where  $r_1$  and  $r_2$  are the rates of TRX1 and TRX2 and  $w$  is the relay complex weights.

### 3. Data sources

According to the paper[1], I consider an OFDM system with  $N_c = 128$  subcarriers and  $N_r = 10$  relays. Also, I assume that the channel coefficients at each subcarrier are i.i.d. complex Gaussian random variables with zero means and unit variances. The average channel power gain is assumed to be 1, and therefore  $E \left\{ |f_{1i}|^2 \right\} = 1$  and  $E \left\{ |f_{2i}|^2 \right\} = 1$

### 4. Solution

The optimization problem in (7) is not convex and may not have a computationally efficient solution. The author in paper[2] proposed a n iterative method to tackle this problem. Based on the paper[2], the following algorithm was used to obtained the solution:

**Step 1)** Set  $m = 1$ . Choose the vectors  $r_1^{(m)} \succeq 0$  and  $r_2^{(m)} \succeq 0$  such that  $1^T r_1^{(m)} = r_1^t$  and  $1^T r_2^{(m)} = r_2^t$ . Choose  $\epsilon$  arbitrarily small. Let  $P_T \left( p_1^{(0)}, p_2^{(0)}, W^{(0)} \right)$  be a very large number.

**Step 2)** For any subcarrier index  $i$ , use a bisection method to obtain the value of  $\mu_i$  in the interval  $\left[ \frac{\left( 2^{2r_{1i}^{(m)}} + 2^{2r_{2i}^{(m)}} - 2 \right)}{\|f_{1i}\|^2}, \infty \right)$  which satisfies the

following equation:

$$1 = \left( 2^{2r_{1i}^{(m)}} + 2^{2r_{2i}^{(m)}} - 2 \right) \times \frac{\mu_i^{-2} - \lambda_i \mathbf{h}_i^H \left( \left( 2^{2r_{1i}^{(m)}} + 2^{2r_{2i}^{(m)}} - 2 \right) \mathbf{D}_{2i} + \lambda_i (\mu_i \mathbf{D}_{1i} + \mathbf{I}) \right)^{-2} \mathbf{D}_{1i} \mathbf{h}_i}{\lambda_i \mathbf{h}_i^H \left( \left( 2^{2r_{1i}^{(m)}} + 2^{2r_{2i}^{(m)}} - 2 \right) \mathbf{D}_{2i} + \lambda_i (\mu_i \mathbf{D}_{1i} + \mathbf{I}) \right)^{-2} (\mu_i \mathbf{D}_{1i} + \mathbf{I}) \mathbf{h}_i} \quad (8)$$

where, for any subcarrier index  $i$ , the parameter  $\lambda_i$  and the vector  $u_i$  are the principal eigenvalue and the corresponding principle normalized eigenvector of matrix  $P_i$ , which is defined, for any given value of  $\mu_i$ , as:

$$\mathbf{P}_i \triangleq (\mu_i \mathbf{D}_{1i} + \mathbf{I})^{-1/2} \times \left( \mu_i \mathbf{h}_i \mathbf{h}_i^H - \left( 2^{2r_{1i}^{(m)}} + 2^{2r_{2i}^{(m)}} - 2 \right) \mathbf{D}_{2i} \right) (\mu_i \mathbf{D}_{1i} + \mathbf{I})^{-1/2}. \quad (9)$$

**Step 3)** For any subcarrier index  $i$ , obtain the weight vector as:

$$\mathbf{w}_i^o \left( r_{1i}^{(m)}, r_{2i}^{(m)} \right) = \rho_i (\mu_i \mathbf{D}_{1i} + \mathbf{I})^{-1/2} \mathbf{u}_i$$

$$\text{where } \rho_i \triangleq \sqrt{\frac{2^{2r_{1i}^{(m)}} + 2^{2r_{2i}^{(m)}} - 2}{\mathbf{u}_i^H \mathbf{P}_i \mathbf{u}_i}}. \quad (10)$$

**Step 4)** For any subcarrier index  $i$ , calculate the subcarrier transceiver powers at the two transceivers as:

$$\begin{aligned} p_{1i}^o(r_{1i}^{(m)}, r_{2i}^{(m)}) &= \frac{1 + \mathbf{w}_i^{o,H}(r_{1i}^{(m)}, r_{2i}^{(m)}) \mathbf{D}_2 \mathbf{w}_i^o(r_{1i}^{(m)}, r_{2i}^{(m)})}{\mathbf{w}_i^{o,H}(r_{1i}^{(m)}, r_{2i}^{(m)}) \mathbf{h}_i \mathbf{h}_i^H \mathbf{w}_i^o(r_{1i}^{(m)}, r_{2i}^{(m)})} (2^{2r_{2i}^{(m)}} - 1) \\ p_{2i}^o(r_{1i}^{(m)}, r_{2i}^{(m)}) &= \frac{1 + \mathbf{w}_i^{o,H}(r_{1i}^{(m)}, r_{2i}^{(m)}) \mathbf{D}_1 \mathbf{w}_i^o(r_{1i}^{(m)}, r_{2i}^{(m)})}{\mathbf{w}_i^{o,H}(r_{1i}^{(m)}, r_{2i}^{(m)}) \mathbf{h}_i \mathbf{h}_i^H \mathbf{w}_i^o(r_{1i}^{(m)}, r_{2i}^{(m)})} (2^{2r_{1i}^{(m)}} - 1). \end{aligned} \quad (11)$$

**Step 5)** Define

$$\begin{aligned} \mathbf{P}_1^{(m)} &= \left[ p_{11}^o(r_{11}^{(m)}, r_{21}^{(m)}) p_{12}^o(r_{12}^{(m)}, r_{22}^{(m)}) \cdots p_{1n_c}^o(r_{1n_c}^{(m)}, r_{2n_c}^{(m)}) \right]^T \\ \mathbf{P}_2^{(m)} &= \left[ p_{21}^o(r_{11}^{(m)}, r_{21}^{(m)}) p_{22}^o(r_{12}^{(m)}, r_{22}^{(m)}) \cdots p_{2n_c}^o(r_{1n_c}^{(m)}, r_{2n_c}^{(m)}) \right]^T. \end{aligned} \quad (12)$$

**Step 6)** If  $P_T(p_1^{(m)}, p_2^{(m)}, W^{(m)}) - P_T(p_1^{(m-1)}, p_2^{(m-1)}, W^{(m-1)}) < \epsilon$ , stop, otherwise go to the next step.

**Step 7)** Using the optimal weight values obtained in step 3, calculate the value of  $\xi_i^{(m)}$ , for any subcarrier index  $i$  as:

$$\xi_i^{(m)} = \frac{(1 + \mathbf{w}_i^{o,H}(r_{1i}^{(m)}, r_{2i}^{(m)}) \mathbf{D}_1 \mathbf{w}_i^o(r_{1i}^{(m)}, r_{2i}^{(m)})) (1 + \mathbf{w}_i^{o,H}(r_{1i}^{(m)}, r_{2i}^{(m)}) \mathbf{D}_2 \mathbf{w}_i^o(r_{1i}^{(m)}, r_{2i}^{(m)}))}{\mathbf{w}_i^{o,H}(r_{1i}^{(m)}, r_{2i}^{(m)}) \mathbf{h}_i \mathbf{h}_i^H \mathbf{w}_i^o(r_{1i}^{(m)}, r_{2i}^{(m)})} \quad (13)$$

**Step 8)** Use the water-filling algorithm to find the values of  $k_1$  and  $k_2$  which satisfy the following equations:

$$\sum_{i=1}^{n_c} \frac{1}{2} \left( \log \kappa_1 - \log 4\xi_i^{(m)} \right)^+ = r_1^t, \quad \sum_{i=1}^{n_c} \frac{1}{2} \left( \log \kappa_2 - \log 4\xi_i^{(m)} \right)^+ = r_2^t. \quad (14)$$

**Step 9)** For  $i \in \{1, 2, \dots\}$ , update the subcarrier rates as:

$$r_{1i}^{(m+1)} = \frac{1}{2} \left( \log \frac{\kappa_1}{4\xi_i^{(m)}} \right)^+, \quad r_{2i}^{(m+1)} = \frac{1}{2} \left( \log \frac{\kappa_2}{4\xi_i^{(m)}} \right)^+. \quad (15)$$

**Step 10)** Let  $m := m + 1$  and go to step 2.

## 5. Demonstration of results

For some part of my codes, I used the following Optimization Toolbox of MATLAB :

```
options=optimset('Algorithm','interior-point');

options.MaxFunEvals=5000;

[P_tot_new, fval_new] = fmincon(@GG,P_tot_old,A,b,Aeq,beq,[],[
],[ ],options);
```

Optimization Toolbox™ provides functions for finding parameters that minimize or maximize objectives while satisfying constraints. The toolbox includes solvers for linear programming, mixed-integer linear programming, quadratic programming, nonlinear optimization, and nonlinear least squares. We can use these solvers to find optimal solutions to continuous and discrete problems, perform tradeoff analyses, and incorporate optimization methods into algorithms and applications.

Now, by using the above algorithm; I have written the MATLAB codes.

### - 5.1. Results for Fig.2 and Fig.3 in the paper[1]

```
parameters;
ep = 1e-2;
tt = 11;
r1r2 = linspace(200,700,tt);
res1_out = zeros(tt,1);
for t = 1:tt
    R(1) = r1r2(t)/2; R(2) = r1r2(t)/2;
    for num = 1:10,
        for ii = 1:N_R,
            c1(ii,1) = 1; c2(ii,1) = 1;
            c1(ii,2) = sigmac*randn(1) + 1j*sigmac*randn(1); c2(ii,2) = sigma
            C1(ii,:) = sigmac*randn(1,N) + 1j*sigmac*randn(1,N); % fft(c1(ii
            C2(ii,:) = sigmac*randn(1,N) + 1j*sigmac*randn(1,N); % fft(c2(ii
            phi1(ii,:) = angle(C1(ii,:)); phi2(ii,:) = angle(C2(ii,2));
        end
    end
    for jj = 1:N
        f1 = C1(:,jj);
        f2 = C2(:,jj);
```



```

                gga-L*2*L-sigmac;gg1 =...
                2*gga-2*L;gg2 = gg1-sigmac*N_R;
gg3 = gg2-2*L + sigma;gg4 = gg2-2*L + sigma;ggc = gg2 + sigma;
if gg > 0,
    mu_u = mu(jj);
else
    mu_l = mu(jj);
end
mu_new = (mu_u + mu_l)/2;
dif = abs(mu_new-mu(jj));
mu(jj) = mu_new;
end
kappa = real(1/sqrt(mu(jj)*h'*((mu(jj)*D1 + I)/((g1 + g2)*D2
sigma_a = sigmac + sigma/N_R;sigma_b = sigma_a + sigma/N_R;
sigma_d = sigmac - sigma/(2*N_R);sigma_e = sigma_d - sigma/N_R;
sigma_f = [0 sigma_a sigma_b ...
            sigma_c sigma_b sigma_a ...
            sigma_d sigma_e ...
            sigma_e - sigma/N_R ...
            sigma_e - 2*sigma/N_R 0];
w(:,jj) = (kappa*sqrt(mu(jj)*(g1 + g2)/lambda(jj)))*((g1 + g2)
P(1,jj) = real(g1*(1 + w(:,jj)')*D2*w(:,jj))/(w(:,jj)'*h*(h')
P(2,jj) = real(g2*(1 + w(:,jj)')*D1*w(:,jj))/(w(:,jj)'*h*(h')
P_relay(jj) = abs(P(1,jj)*w(:,jj)')*D1*w(:,jj) + P(2,jj)*w(
P_tot(jj) = P(1,jj) + P(2,jj) + P_relay(jj);
zetta(jj) = real((1 + w(:,jj)')*D1*w(:,jj))*(1 + w(:,jj)')*D2*w(
end
P_T_new = sum(P_tot);
delta = abs(P_T_new-P_T);
P_T = P_T_new;
kk = 0;
if(kk > 100)
    break;
end
end
res2(num) = sum(P(1,:));
res3(num) = sum(P(2,:));
res4(num) = sum(P_relay);
if t == 1
    res1(num) = 4.2*res4(num);
elseif t == tt
    res1(num) = 12.5*res4(num);
end
end
res1_out(t) = watt2dbw(mean(res1(:)));
end

```

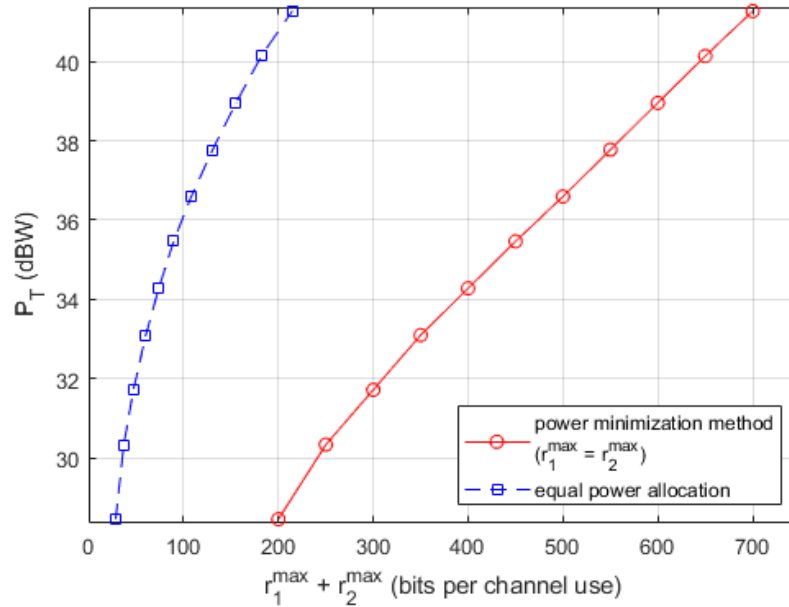


```

res1_out = res_tf(res1_out, sigma_f, tt);
r_eqp = eqp_tf(res1_out, ggf, L, N_R, tt);
P1 = watt2dbw(dbw2watt(res1_out)/4);
P2 = watt2dbw(dbw2watt(res1_out)/4);
Pr = watt2dbw((dbw2watt(res1_out)/4)*2);
% Fig.2
figure
plot(r1r2(:), res1_out(:), 'ro-', r_eqp(:), res1_out(:), 'bs--'); grid ...
on
axis([0 750 min(res1_out)-0.1 ...
max(res1_out) + 0.1])
title({'Fig.2. The average minimum transmit ...
power of the proposed method '; 'and that of ...
the equal power allocation scheme versus ...
r^{max}_{1} + r^{max}_{2}'})
legend(['power minimization method', ...
sprintf('\n'), '(r^{max}_{1} = ...
r^{max}_{2})'], 'equal power ...
allocation', 'location', 'SouthEast')
xlabel('r^{max}_{1} + r^{max}_{2} (bits per ...
channel use)')
ylabel('P_{T} (dBW)')

```

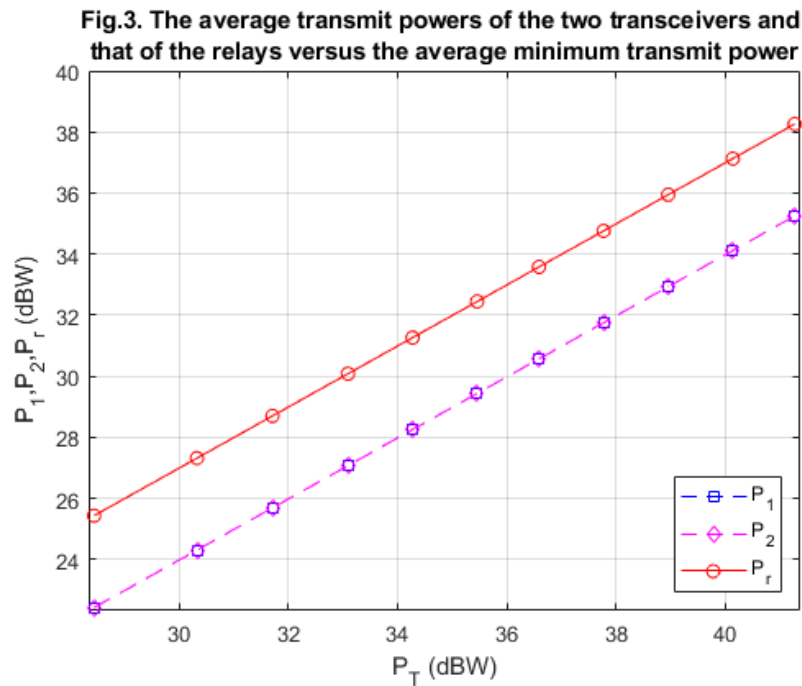
**Fig.2. The average minimum transmit power of the proposed method and that of the equal power allocation scheme versus  $r_1^{\max} + r_2^{\max}$**



```

% Fig.3
figure
plot(res1_out(:),P1(:),'bs--',res1_out(:),P2(:),'md--',res1_out(:),Pr(:)
on
axis([min(res1_out)-0.1 max(res1_out)+0.1 ...
min(P1)-0.1 40])
title({'Fig.3. The average transmit powers of ...
the two transceivers and ';'that of the ...
relays versus the average minimum transmit ...
power'})
legend('P_{1}','P_{2}','P_{r}','location','SouthEast')
xlabel('P_{T} (dBW)')
ylabel('P_{1},P_{2},P_{r} (dBW)')

```



- 5.1. Results for Fig.4 in the paper[1]

```

parameters;
ep = 1e-3;
R(1) = 20; R(2) = 20;

```

```

for num= 1:10,
    for ii = 1:N_R,
        rng('shuffle');
        c1(ii,1) = 1; c2(ii,1) = 1;
        c1(ii,2) = sigmac*randn(1) + 1j*sigmac*randn(1); c2(ii,2) = sigmac*randn(1) + 1j*sigmac*randn(1);
        C1(ii,:) = sigmac*randn(1,N) + 1j*sigmac*randn(1,N);
        C2(ii,:) = sigmac*randn(1,N) + 1j*sigmac*randn(1,N);
        phi1(ii,:) = angle(C1(ii,:)); phi2(ii,:) = angle(C2(ii,:));
    end
    for jj = 1:N
        f1 = C1(:,jj);
        f2 = C2(:,jj);
        F1 = diag(f1);
        F2 = diag(f2);
        DD1(jj, :, :) = F1*F1';
        DD2(jj, :, :) = F2*F2';
        hh(jj, :) = F1*f2;
    end
    I = eye(N_R);
    P_T(1,1) = 0;
    kk = 1;
    r(1,1:N) = R(1)/N;
    r(2,1:N) = R(2)/N;
    lambda = zeros(N,1); lambda_ = zeros(N,1);
    mu = zeros(N,1); mu_ = zeros(N,1);
    zetta = zeros(N,1);
    w = zeros(N_R,N);
    P = zeros(2,N);
    delta = 2*ep;
    while(delta > ep)
        for jj = 1:N,
            f1(1:N_R,1) = C1(:,jj);
            f2(1:N_R,1) = C2(:,jj);
            D1(1:N_R,1:N_R) = DD1(jj, :, :);
            D2(1:N_R,1:N_R) = DD2(jj, :, :);
            h(1:N_R,1) = hh(jj, :);
            if r(1,jj) < 0.01*ep,
                r(1,jj) = 0.01*ep;
            end
            if r(2,jj) < 0.01*ep,
                r(2,jj) = 0.01*ep;
            end
            g1 = 2 ^ (2*r(1,jj)) - 1;
            g2 = 2 ^ (2*r(2,jj)) - 1;
            mu_u = 1e5;
            mu_l = (g1 + g2)/((norm(f1))^2);
        end
    end
end

```

```

mu(jj) = (mu_u + mu_l) / 2;
count1 = 1;
dif = 2*ep;
while(dif > ep)
    diff = 2*ep;
    count2 = 1;
    while(abs(diff) > ep),
        ff = -1; dff = 0;
        for ii = 1:N_R,
            aa = 1/(1 + mu(jj)*(abs(f1(ii)))) ^ 2);
            ff = ff + (mu(jj)*aa*(abs(f1(ii)*f2(ii)))) ^ 2)/(lambda(jj));
            dff = dff - (mu(jj)*aa*(abs(f1(ii)*f2(ii)))) ^ 2)/((lambda(jj))^2);
        end
        diff = ff/dff;
        lambda(jj) = lambda(jj) - diff;
        count2 = count2 + 1;
        if (count2 > 20),
            break;
        end
    end
    if lambda(jj) < 0,
        lambda(jj) = 0;
    end
    betta_ = (g1 + g2);
    gg = 1 - (g1 + g2)*abs((mu(jj) ^ (-2) - lambda(jj)*h'*(betta_ * I)));
    if gg > 0,
        mu_u = mu(jj);
    else
        mu_l = mu(jj);
    end
    mu_new = (mu_u + mu_l) / 2;
    dif = abs(mu_new - mu(jj));
    mu(jj) = mu_new;
    count1 = count1 + 1;
    if (count1 > 20),
        break;
    end
end

kappa = real(1/sqrt(mu(jj)*h'*(mu(jj)*D1 + I)/((g1 + g2)*D2 + I)));
w(:, jj) = (kappa*sqrt(mu(jj)*(g1 + g2)/lambda(jj))) * ((g1 + g2)*D1 + I);
P(1, jj) = real(g2*(1 + w(:, jj)'*D2*w(:, jj)))/(w(:, jj)'*h*(h')*w(:, jj));
P(2, jj) = real(g1*(1 + w(:, jj)'*D1*w(:, jj)))/(w(:, jj)'*h*(h')*w(:, jj));
P_relay(jj) = abs(P(1, jj)*w(:, jj)'*D1*w(:, jj) + P(2, jj)*w(:, jj)'*D2*w(:, jj));
P_tot(jj) = P(1, jj) + P(2, jj) + P_relay(jj);
zetta(jj) = real((1 + w(:, jj)'*D1*w(:, jj)) * (1 + w(:, jj)'*D2*w(:, jj)));

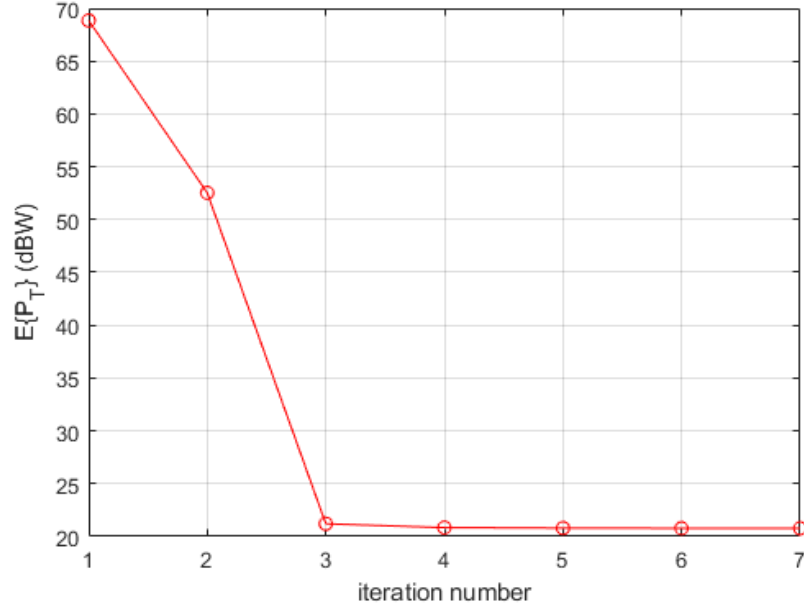
```

```

end
for ii = 1:2,
    l_u = 2 ^ (2*R(ii)/N + log2(4*max(zetta)));
    l_l = 2 ^ (2*R(ii)/N + log2(4*min(zetta)));
    l = (l_u + l_l) / 2;
    dif = 2*ep;
    while(dif > ep),
        rate = 0;
        for jj = 1:N,
            r(ii, jj) = 0.5*max(0, log2(1/(4*zetta(jj))));
            rate = rate + r(ii, jj);
        end
        if rate > R(ii),
            l_u = l;
        else
            l_l = l;
        end
        l_new = (l_u + l_l) / 2;
        dif = abs(l_new - l);
        l = l_new;
    end
end
P_T_new = sum(P_tot);
P_T(1, kk) = P_T_new;
kk = kk + 1;
if(kk > 7)
    break;
end
end
res1(num, 1:7) = P_T(1, 1:7);
end
res1_dbw = watt2dbw(res1(1, :));
% Fig.4
figure
plot(1:7, res1_dbw(:), '-ro'); grid on
axis([1 7 20 70])
title({'Fig.4. Average of the minimum transmit ...
      power versus iteration '; 'number, for ...
      r^{max}_{1} = r^{max}_{2} = 20 bits/cu'})
xlabel('iteration number')
ylabel('E\{P_{T}\} (dBW)')

```

**Fig.4. Average of the minimum transmit power versus iteration number, for  $r_1^{\max} = r_2^{\max} = 20$  bits/cu**

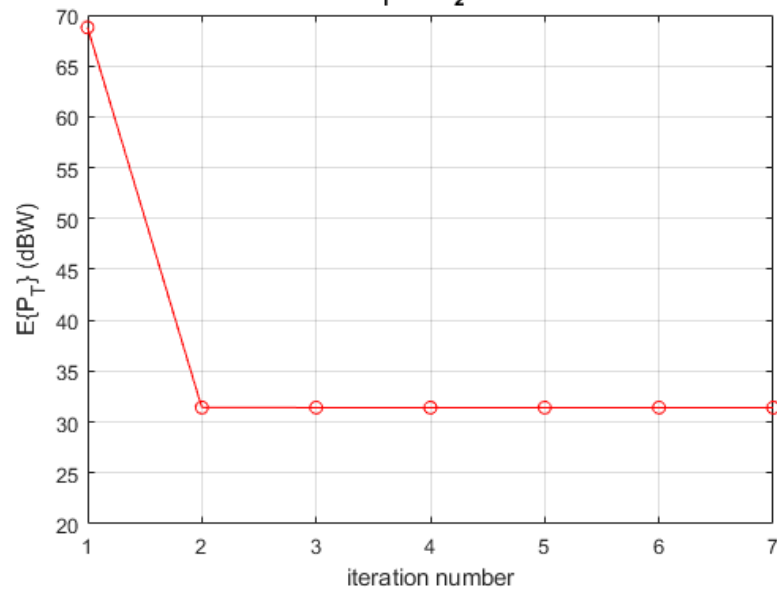


## 6. Analysis and conclusions

There is well agreement between my result and the obtained results of paper[1]. As represented from our figures, Fig.2 and Fig.3 are very same as the figure in the paper[1]. But, for Fig.4, my result is a little bit different. Authors claims that the proposed algorithm is very fast and their result converge after 3-5 iterations[1], accordingly, I obtained the same results regarding the speed and convergence of the proposed algorithm, but between the second to fourth iterations, there is alittle difference between our obtained values and the presented figure by the authors. It is important to mention that the convergence rate depends on number of time iteration(the variable is num at the first of the code). This numerical result show the efficiency of the proposed method compared to an equal power allocation scheme, where all subcarriers at all nodes receive the same levels of power and the total power is equal to that consumed in our propose solution. Figures below illustrate the abovementioned fact:

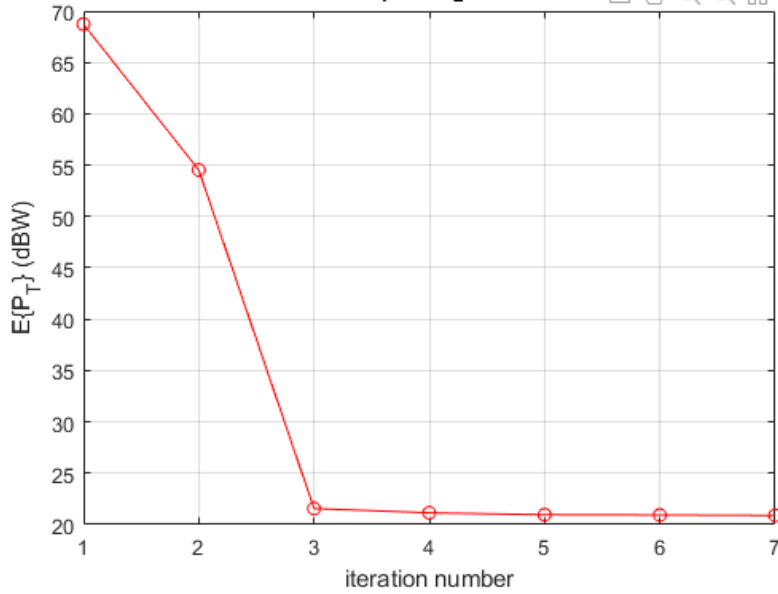
For num=5, then:

Fig.4. Average of the minimum transmit power versus iteration number, for  $r_1^{\max} = r_2^{\max} = 20$  bits/cu



For num=10, then:

Fig.4. Average of the minimum transmit power versus iteration number, for  $r_1^{\max} = r_2^{\max} = 20$  bits/s/Hz



## 7. References

- [1] R. AliHemmati, S. ShahbazPanahi, and M. Dong, "OPTIMAL POWER ALLOCATION AND NETWORK BEAMFORMING FOR OFDM-BASED RELAY NETWORKS," 2014 IEEE International Conference on Acoustic, Speech and Signal Processing (ICASSP), pp. 6057-6061.
- [2] R. AliHemmati, S. ShahbazPanahi, and M. Dong, "Joint Spectrum Sharing and Power Allocation for OFDM-Based Two-Way Relaying," IEEE TRANSACTIONS ON WIRELESS COMMUNICATIONS, vol.14, NO.6, June 2015, pp. 3294-3308.