

ECE 602 – Section 6
Lagrangian Duality

- Lagrangian and Lagrangian multipliers
- Dual variables and dual function
- Weak vs strong duality
- Saddle-point theorem
- Karush-Kuhn-Tucker conditions
- Alternating Direction Method of Multipliers

- Consider the problem

$$\begin{aligned} \min_x \quad & f_0(x) \\ \text{s.t.} \quad & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p \end{aligned}$$

with $\mathcal{D} = \bigcap_{i=0}^m \mathbf{dom} f_i \cap \bigcap_{i=1}^p \mathbf{dom} h_i \neq \emptyset$ and optimal value p^* .

- We define the associated *Lagrangian function* as

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

with the variables $x \in \mathbf{R}^n$, $\lambda \in \mathbf{R}_+^m$ and $\nu \in \mathbf{R}^p$.

- The vectors λ and ν are called the *dual variables* or *Lagrange multipliers*. In this case, x is referred to as the *primal* variable.

- The *Lagrange dual function* is defined as

$$g(\lambda, \nu) = \inf_{x \in \mathcal{D}} L(x, \lambda, \nu)$$

- When the Lagrangian is unbounded below in x , the dual function takes on the value $-\infty$.
- The *dual function is always concave*, as its cost is a point-wise infimum over affine functions.
- Moreover, one can shown that, for any $\lambda \succeq 0$ and ν , we have

$$g(\lambda, \nu) \leq p^*$$

- We refer to (λ, ν) , with $\lambda \succeq 0$ and $(\lambda, \nu) \in \mathbf{dom} g$, as *dual feasible*.

- Consider the problem

$$\begin{aligned} \min_x \quad & x^T x \\ \text{s.t.} \quad & Ax = b. \end{aligned}$$

with its associated Lagrangian $L(x, \nu) = x^T x + \nu^T (Ax - b)$.

- To minimize $L(x, \nu)$ over x we solve $\nabla_x L(x, \nu) = 2x + A^T \nu = 0$, which yields $x = -(1/2)A^T \nu$. Then, substituting the latter into $L(x, \nu)$ leads to

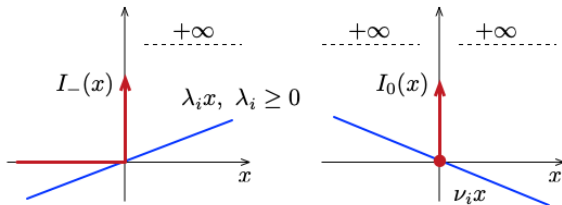
$$g(\nu) = L(- (1/2)A^T \nu, \nu) = -(1/4)\nu^T AA^T \nu - b^T \nu,$$

- The dual $g(\nu)$ is obviously a concave function. Moreover, by the *low-bound property*, we have

$$-(1/4)\nu^T AA^T \nu - b^T \nu \leq p^*$$

- Let $I_-(x)$ and $I_0(x)$ be the indicator functions of the sets $-\mathbf{R}_+^n$ and $\{0\}$, respectively. Then, the original (*primal*) problem can be expressed in an *unconstrained form* as

$$\min_x f_0(x) + \sum_{i=1}^m \underbrace{I_-(f_i(x))}_{\approx \lambda_i f_i(x)} + \sum_{i=1}^p \underbrace{I_0(h_i(x))}_{\approx \nu_i h_i(x)}$$



- The idea behind the definition of L is to *approximate* (“soften”) $I_-(x)$ and $I_0(x)$ by their *linear under-estimators*.
- Note that: (1) the approximations are *exact* at 0, and (2) the “slopes” are *variables* (and hence adjustable).

- Consider the following *linearly constrained* problem

$$\min_x f_0(x) \quad \text{s.t.} \quad Ax \preceq b, \quad Cx = d$$

- In this case, the dual function can be expressed as

$$\begin{aligned} g(\lambda, \nu) &= \inf_x (f_0(x) + \lambda^T (Ax - b) + \nu^T (Cx - d)) = \\ &= -b^T \lambda - d^T \nu + \inf_x (f_0(x) + (A^T \lambda + C^T \nu)^T x) = \\ &= -b^T \lambda - d^T \nu - \underbrace{\sup_x ((-A^T \lambda - C^T \nu)^T x - f_0(x))}_{f_0^*(-A^T \lambda - C^T \nu)} \end{aligned}$$

- Therefore, we have

$$g(\lambda, \nu) = -b^T \lambda - d^T \nu - f_0^*(-A^T \lambda - C^T \nu)$$

with $\mathbf{dom} \, g = \{(\lambda, \nu) \mid -A^T \lambda - C^T \nu \in \mathbf{dom} \, f_0^*\}$.

- Consider the following problem

$$\begin{aligned} \min_x \quad & \|x\| \\ \text{s.t.} \quad & Ax = b \end{aligned}$$

where $\|\cdot\|$ is any norm.

- The conjugate of $\|x\|$ is given by the indicator function of the unit ball $\mathcal{B}^* = \{u \mid \|u\|_* \leq 1\}$, namely

$$f_0^*(y) = I_{\mathcal{B}^*}(y) = \begin{cases} 0 & \|y\|_* \leq 1 \\ \infty & \text{otherwise,} \end{cases}$$

- Consequently, we have

$$g(\nu) = -b^T \nu - I_{\mathcal{B}^*}(-A^T \nu)$$

- Consider an LP in standard form

$$\begin{aligned} \min_x \quad & c^T x \\ \text{s.t.} \quad & Ax = b, \quad x \succeq 0 \end{aligned}$$

- The associated Lagrangian is given by

$$L(x, \lambda, \nu) = c^T x - \lambda^T x + \nu^T (Ax - b) = -\nu^T b + (c + A^T \nu - \lambda)^T x$$

- The dual function is

$$\begin{aligned} g(\lambda, \nu) &= \inf_x L(x, \lambda, \nu) = -\nu^T b + \inf_x (c + A^T \nu - \lambda)^T x \\ &= \begin{cases} -\nu^T b, & \text{if } c + A^T \nu - \lambda = 0 \\ -\infty, & \text{otherwise} \end{cases} \end{aligned}$$

- The lower bound property is nontrivial only when $A^T \nu - \lambda + c = 0$ and $\lambda \succeq 0$. When this occurs, $-\nu^T b$ is a *lower bound* on the optimal value of the LP.

- For each pair (λ, ν) with $\lambda \succeq 0$, the Lagrange dual function gives us a lower bound on p^* .
- The *best lower bound* can be found by solving

$$\begin{aligned} \max_{\lambda, \nu} & g(\lambda, \nu) \\ \text{s.t.} & \lambda \succeq 0, \end{aligned}$$

which is called the *Lagrange dual problem*.

- We refer to (λ^*, ν^*) as the *dual optimal* or *optimal Lagrange multipliers*.
- The dual problem is always a convex optimization problem, *regardless* whether or not f_0 is convex.

EXAMPLE: LAGRANGE DUAL OF STANDARD FORM LP

- Consider the following LP problem in *standard form*

$$\begin{aligned} \min_x \quad & c^T x \\ \text{s.t.} \quad & Ax = b, \quad x \succeq 0 \end{aligned}$$

- Its Lagrange dual function is given by

$$g(\nu) = \begin{cases} -b^T \nu, & A^T \nu - \lambda + c = 0 \\ -\infty, & \text{otherwise} \end{cases}$$

- Therefore, the Lagrange dual problem of the standard form LP is

$$\begin{aligned} \max_x \quad & -b^T \nu \\ \text{s.t.} \quad & A^T \nu - \lambda + c = 0, \quad \lambda \succeq 0 \end{aligned}$$

or, equivalently,

$$\begin{aligned} \max_x \quad & -b^T \nu \\ \text{s.t.} \quad & A^T \nu + c \succeq 0 \end{aligned}$$

which is an LP in *inequality form*

- Consider the following LP problem in *inequality form*

$$\begin{aligned} \min_x \quad & c^T x \\ \text{s.t.} \quad & Ax \preceq b \end{aligned}$$

- Its Lagrange dual function is given by

$$g(\nu) = \begin{cases} -b^T \lambda, & A^T \lambda + c = 0 \\ -\infty, & \text{otherwise.} \end{cases}$$

Note that λ *dual feasible* if $\lambda \succeq 0$ and $A^T \lambda + c = 0$.

- Therefore, the dual form of the LP in inequality form is

$$\begin{aligned} \max_x \quad & -b^T \lambda \\ \text{s.t.} \quad & A^T \lambda = -c, \lambda \succeq 0 \end{aligned}$$

which is an LP in *standard form*.

- The optimal value of the Lagrange dual problem d^* satisfies

$$d^* \leq p^*$$

This property is called *weak duality*.

- Weak duality holds for both convex and non-convex problems.
 - The difference $p^* - d^*$ is referred to as the *optimal duality gap*.
 - The bound can be used to find lower bounds for difficult problems.
- We say that *strong duality* holds if

$$d^* = p^*$$

- Strong duality does not hold, in general.
- If the *primal* problem is convex, we usually have strong duality.
- The conditions which guarantee strong duality in convex problems are called *constraint qualifications*.

- Consider the optimization problem

$$\begin{aligned}
 \text{(P0)} \quad & \min_x f_0(x) \\
 & \text{s.t. } Ax = b, f_i(x) \leq 0, i = 1, \dots, m
 \end{aligned}$$

- The constraint i is called *active* at point x if $f_i(x) = 0$. Otherwise, if $f_i(x) < 0$, the constraint i is called *inactive* at x .
- The problem **P0** is called *strictly feasible* if there exists $x \in \text{int } \mathcal{D}$ such that

$$Ax = b \quad \text{and} \quad f_i(x) < 0, i = 1, \dots, m$$

or, $f_i(x) \leq 0$, when $f_i(x) = a_i^T x - b_i$. In other words, *all* the inequality constraints are *inactive* at x .

Slater's constraint qualification

If the problem **P0** is *strictly feasible* with an optimal value $p^* > -\infty$, then the dual optimum is attained, implying there exists dual feasible (λ^*, ν^*) such that

$$g(\lambda^*, \nu^*) = d^* = p^*$$

- The optimal values of the primal and dual problems are given by

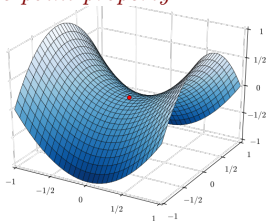
$$p^* = \inf_x \sup_{\lambda \geq 0} L(x, \lambda) \quad \text{and} \quad d^* = \sup_{\lambda \geq 0} \inf_x L(x, \lambda)$$

- Weak duality* can be expressed as the *min-max inequality*

$$\sup_{\lambda \geq 0} \inf_x L(x, \lambda) \leq \inf_x \sup_{\lambda \geq 0} L(x, \lambda)$$

- Strong duality* is characterized by the *saddle-point property*

$$\begin{aligned} \sup_{\lambda \geq 0} \inf_x L(x, \lambda) &= \\ &= \inf_x \sup_{\lambda \geq 0} L(x, \lambda) \end{aligned}$$



- The saddle-point property allows *switching the order* of the “min over x ” and the “max over λ ” without affecting the final result.

- If strong duality holds and a dual optimal solution (λ^*, ν^*) exists, then any primal optimal point is also a minimizer of $L(x, \lambda^*, \nu^*)$.
- This sometimes allows us to compute a *primal* optimal solution from a *dual* optimal solution (that is obtained via convex optimization).
- More precisely, suppose we have strong duality, (λ^*, ν^*) are known and the minimizer of $L(x, \lambda^*, \nu^*)$ is unique. Then, the solution to

$$x^* = \arg \min_x L(x, \lambda^*, \nu^*)$$

with $L(x, \lambda^*, \nu^*) = f_0(x) + \sum_{i=1}^m \lambda_i^* f_i(x) + \sum_{i=1}^p \nu_i^* h_i(x)$, must be primal feasible, and hence primal optimal.

- In practice, $L(x, \lambda^*, \nu^*)$ can be minimized by means of any method of *unconstrained* optimization.

EXAMPLE: ENTROPY MAXIMIZATION

- Consider the *entropy maximization problem*

$$\begin{aligned} \min_x f_0(x) &= \sum_{i=1}^n x_i \log x_i \\ \text{s.t. } Ax &\preceq b, \mathbf{1}^T x = 1 \end{aligned}$$

with $\text{dom } f_0 = \mathbf{R}_{++}^n$. The (convex) set defined by $\mathbf{1}^T x = 1$ is called the *probability simplex*.

- The conjugate of the *negative entropy function* $u \log u$ (with $u > 0$) is $\exp(v - 1)$. Then, due to the separability of f_0 , we have

$$f_0^*(y) = \sum_{i=1}^n \exp(y_i - 1), \quad \text{with } \text{dom } f_0^* = \mathbf{R}^n$$

- Therefore, the associated dual function is given by

$$g(\lambda, \nu) = -b^T \lambda - \nu - \sum_{i=1}^n \exp(-a_i^T \lambda - \nu - 1)$$

where a_i is the i th column of A .

- The dual problem can then be expressed as

$$\begin{aligned} \max_{\lambda, \nu} \quad & -b^T \lambda - \nu - e^{-\nu-1} (\mathbf{1}^T e^{-A^T \lambda}) \\ \text{s.t.} \quad & \lambda \succeq 0 \end{aligned}$$

and Slater condition tells us that the optimal duality gap is zero, if there exists an $x \succ 0$ with $Ax \preceq b$ and $\mathbf{1}^T x = 1$.

- For fixed λ , the objective function is maximized when its derivative w.r.t. ν is zero, i.e.,

$$v = \log (\mathbf{1}^T e^{-A^T \lambda}) - 1$$

- Substituting the above value back into the objective yields

$$\begin{aligned} \max_{\lambda} \quad & -b^T \lambda - \log (\mathbf{1}^T e^{-A^T \lambda}) \\ \text{s.t.} \quad & \lambda \succeq 0 \end{aligned}$$

- The dual problem has the form of a *geometric program* with non-negativity constraints.

- For *any* optimization problem with differentiable f_i and h_i , *for which strong duality holds*, any pair of primal and dual optimal points must satisfy the *Karush-Kuhn-Tucker (KKT) conditions*.
- The KKT conditions characterize the primal optimal x^* and dual optimal (λ^*, ν^*) solutions using four principal statements.
- **Primal feasibility**

$$f_i(x^*) \leq 0, \quad i = 1, \dots, m, \quad h_i(x^*) = 0, \quad i = 1, \dots, p$$

- **Dual feasibility**

$$\lambda_i^* \geq 0, \quad i = 1, \dots, m$$

- **First-order optimality**

$$\nabla f_0(x^*) + \sum_{i=1}^m \lambda_i^* \nabla f_i(x^*) + \sum_{i=1}^p \nu_i^* \nabla h_i(x^*) = 0$$

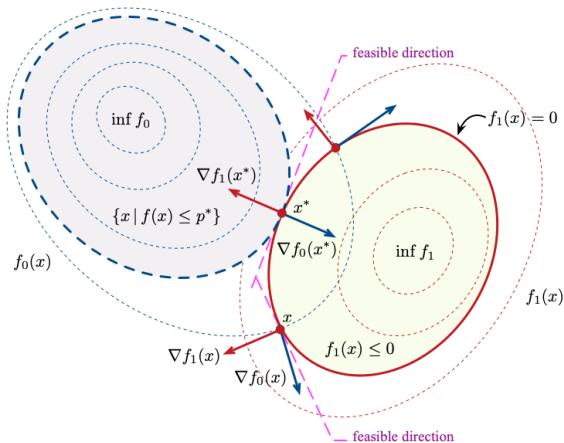
- **Complementary slackness**

$$\lambda_i^* f_i(x^*) = 0, \quad i = 1, \dots, m$$

FIRST-ORDER OPTIMALITY

- To understand “how Lagrange multipliers work”, consider the convex (differentiable) problem

$$\min_x f_0(x) \quad \text{s.t.} \quad f_1(x) \leq 0$$



Can you see the duality?

- The optimal solution is located on the *boundary* of the feasible set.
- At point x , the direction orthogonal to $\nabla f_1(x)$ defines a *feasible direction*. If we move *locally* near x along this tangent direction, the (active) constraint $f_1(x) = 0$ holds.
- However, $\nabla f_0(x)$ is *not* orthogonal to the feasible direction, meaning that, when we move near x along the feasible direction, the values of $f_0(x)$ keep changing.
- At the optimal point x^* , $\nabla f_0(x^*)$ is orthogonal to the feasible direction and, therefore, when we move near x along this line, the value of $f_0(x)$ remains constant.
- In this case, $\nabla f_0(x^*)$ and $\nabla f_1(x^*)$ are *collinear*, implying there exists a *proportionality coefficient* $\lambda^* > 0$ such that

$$\nabla f_0(x^*) + \lambda^* \nabla f_1(x^*) = 0$$

This λ^* is the optimal dual variable.

- Recall that, for $\min_{x: Ax=b} f_0(x)$, we have $\nabla f_0(x^*) + A^T \nu^* = 0$.

- The complimentary slackness condition states that, if $f_i(x^*) < 0$ (i.e., the i th constraint is *inactive*), $\lambda_i^* = 0$. At the same time, if $f_i(x^*) = 0$ (i.e., the constraint is *active*), $\lambda_i^* > 0$.
- Recall that, due to the strong duality, we have $f_0(x^*) = L(x^*, \lambda^*, \nu^*)$. Since the inactive inequality constraints could potentially disturb this “balance”, their constraint functions are multiplied by $\lambda_i^* = 0$.
- For the active constraints, the “slopes” $\lambda_i^* > 0$ are set according to the first-order optimality condition. However, the factual values of λ_i^* have no effect on the optimal value, since they are multiplied by $f_i(x^*) = 0$.
- The inactive constraints are, in fact, *redundant*.

- Consider the equality constrained problem

$$\min_x f_0(x) \quad \text{s.t.} \quad Ax = b$$

for some $A \in \mathbf{R}^{m \times n}$ and $b \in \mathbf{R}^m$.

- For this problem, the *augmented Lagrangian function* is defined as

$$L_\rho(x, \nu) = f_0(x) + \nu^T (Ax - b) + \frac{\rho}{2} \|Ax - b\|_2^2$$

for some *penalty parameter* $\rho > 0$. Note that, at the optimal solution, $L(x^*, \nu^*) = L_\rho(x^*, \nu^*) = f_0(x^*)$.

- The benefit of including the penalty term is that $g_\rho(\nu) = \inf_x L_\rho(x, \nu)$ can be shown to be differentiable under rather mild conditions on f_0 .
- By completing the squares, the augmented Lagrangian function can also be expressed as

$$L_\rho(x, \nu) = f_0(x) + \frac{\rho}{2} \|Ax - b + \nu/\rho\|_2^2 - \frac{1}{2\rho} \|\nu\|_2^2$$

- The *Method of Multiplier* iterates between the estimates of x and ν according to

$$x^{(t+1)} = \arg \min_x L_\rho(x, \nu^{(t)})$$

$$\nu^{(t+1)} = \nu^{(t)} + \rho(Ax^{(t+1)} - b)$$

- The first step *minimizes* L_ρ w.r.t. the primal variable, with the dual variable being fixed at its "old" value $\nu^{(t)}$.
- The second step *maximizes* L_ρ w.r.t. the dual variable by means of *gradient ascent* with a step-size ρ .
- Note that the feasibility requires $Ax^* - b = 0$ and $\nabla f_0(x^*) + A^T \nu^* = 0$. Since $x^{(t+1)}$ minimizes $L_\rho(x, \nu^{(t)})$, we have

$$0 = \nabla_x L_\rho(x^{(t+1)}, \nu^{(t)}) = \nabla f_0(x^{(t+1)}) + A^T \underbrace{(\nu^{(t)} + \rho(Ax^{(t+1)} - b))}_{\nu^{(t+1)}}$$

- By using ρ as the step size, the iterate $(x^{(t+1)}, \nu^{(t+1)})$ is *dual feasible*.

- Consider the following *unconstrained* problem

$$\min_x f(Ax) + g(Bx)$$

with $A \in \mathbf{R}^{m \times n}$, $B \in \mathbf{R}^{p \times n}$ and some sub-differentiable f and g .

- Using two auxiliary variables u and v , the problem can be expressed in a *constrained form* as

$$\begin{aligned} \min_{x,u,v} f(u) + g(v) \\ \text{s.t. } Ax = u, Bx = v \end{aligned}$$

- The associated augmented Lagrangian is given by

$$\begin{aligned} L_{\rho_u, \rho_v}(x, u, v, \nu_u, \nu_v) = f(u) + g(v) + \\ + \frac{\rho_u}{2} \underbrace{\|Ax - u + \nu_u / \rho_u\|_2^2}_{p_u} + \frac{\rho_v}{2} \underbrace{\|Bx - v + \nu_v / \rho_v\|_2^2}_{p_v} + \text{resid.} \end{aligned}$$

where p_u and p_v can be viewed as *scaled* Lagrange multipliers.

- The *Alternating Direction Method of Multipliers* (ADMM) updates the primal variable *one at a time*, followed by gradient ascent on the dual variables. The ADMM consists of a *primal* and a *dual* update.

Primal Update

$$\begin{aligned} x^{(t+1)} &= \arg \min_x \left\{ \frac{\rho_u}{2} \|Ax - u^{(t)} + p_u^{(t)}\|_2^2 + \frac{\rho_v}{2} \|Bx - v^{(t)} + p_v^{(t)}\|_2^2 \right\} = \\ &= (\rho_u A^T A + \rho_v B^T B)^{-1} (\rho_u A^T (u^{(t)} - p_u^{(t)}) + \rho_v B^T (v^{(t)} - p_v^{(t)})) \end{aligned}$$

$$\begin{aligned} u^{(t+1)} &= \arg \min_u \left\{ f(u) + \frac{\rho_u}{2} \|Ax^{(t+1)} - u + p_u^{(t)}\|_2^2 \right\} = \\ &= \mathbf{prox}_{(1/\rho_u)f} (Ax^{(t+1)} + p_u^{(t)}) \end{aligned}$$

$$\begin{aligned} v^{(t+1)} &= \arg \min_v \left\{ g(v) + \frac{\rho_v}{2} \|Bx^{(t+1)} - v + p_v^{(t)}\|_2^2 \right\} = \\ &= \mathbf{prox}_{(1/\rho_v)g} (Bx^{(t+1)} + p_v^{(t)}) \end{aligned}$$

Dual Update

$$p_u^{(t+1)} = p_u^{(t)} + Ax^{(t+1)} - u^{(t+1)}$$

$$p_v^{(t+1)} = p_v^{(t)} + Bx^{(t+1)} - v^{(t+1)}$$

- If both f and g are convex, it can be shown that $f(u^{(t)}) + g(v^{(t)}) \rightarrow p^*$ as $t \rightarrow \infty$, i.e., the objective function of the iterates approaches p^* .
- It is often the case that ADMM converges to *modest accuracy* – sufficient for many applications – within a few tens of iterations.
- The method is not symmetric in the order of updates.
- In theory, the ADMM is guaranteed to converge for any values of the penalty parameters ρ_u and ρ_v (e.g., $\rho_u = \rho_v = 1$). In practice, however, their choice can influence the rate of convergence.
- There are several heuristics for setting ρ_u and ρ_v *adaptively*.

EXAMPLE: PARALLEL PROJECTIONS

- Let $\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_N$ be a set of N convex and closed sets in \mathbf{R}^n . Consider the problem of finding a point in the intersection $\bigcap_{i=1}^N \mathcal{A}_i$.

- Let \mathcal{C} and \mathcal{D} be two subsets in \mathbf{R}^{nN} defined as

$$\mathcal{C} = \mathcal{A}_1 \times \mathcal{A}_2 \times \dots \times \mathcal{A}_N$$

$$\mathcal{D} = \{(x_1, x_2, \dots, x_N) \in \mathbf{R}^{nN} \mid x_1 = x_2 = \dots = x_N\}$$

- If $(x_1, x_2, \dots, x_N) \in \mathbf{R}^{nN}$, then its *orthogonal projection* onto \mathcal{C} is given by

$$\mathcal{P}_{\mathcal{C}}(x) = \left(\mathcal{P}_{\mathcal{A}_1}(x_1), \mathcal{P}_{\mathcal{A}_2}(x_2), \dots, \mathcal{P}_{\mathcal{A}_N}(x_N) \right) \in \mathcal{C} \subset \mathbf{R}^{nN}$$

- Let $\bar{x} = (1/N) \sum_{i=1}^N x_i$ be the *average* of x_1, x_2, \dots, x_N . Then, the *orthogonal projection* onto \mathcal{D} is given by

$$\mathcal{P}_{\mathcal{D}}(x) = \left(\overbrace{\bar{x}, \bar{x}, \dots, \bar{x}}^{n \text{ times}} \right) \in \mathcal{D} \subset \mathbf{R}^{nN}$$

- Then, the ADMM iterations are defined as

$$x_i^{(t+1)} = \mathcal{P}_{\mathcal{A}_i}(u^{(t)} - p_i^{(t)}), \forall i$$

$$u^{(t+1)} = \frac{1}{N} \sum_{i=1}^N (x_i^{(t+1)} + p_i^{(t)})$$

$$p_i^{(t+1)} = p_i^{(t)} + x_i^{(t+1)} - u^{(t+1)}, \forall i$$

with $u^{(t)} \in \mathbf{R}^n$ and $(p_1^{(t)}, p_2^{(t)}, \dots, p_N^{(t)}) \in \mathbf{R}^{nN}$.

- The iterations can also be redefined in a simplified form as

$$x_i^{(t+1)} = \mathcal{P}_{\mathcal{A}_i}(\bar{x}^{(t)} - p_i^{(t)}), \forall i$$

$$p_i^{(t+1)} = p_i^{(t)} + (x_i^{(t+1)} - \bar{x}^{(t+1)}), \forall i$$

- Note that $p_i^{(t)}$ is the running sum of the “discrepancies” $x_i^{(t)} - \bar{x}^{(t)}$ (if we assume $p_i^{(0)} = 0$, for all i).

EXAMPLE: QUADRATIC PROGRAMMING

- Recall that the *standard form* QP is given by

$$\begin{aligned} \min_x \quad & (1/2) x^T P x + q^T x \\ \text{s.t.} \quad & A x = b, \quad x \succeq 0 \end{aligned}$$

for some $P \in \mathbf{S}_{++}^n$, $A \in \mathbf{R}^{m \times n}$, $q \in \mathbf{R}^n$ and $b \in \mathbf{R}^m$. Note that, if P is zero, the QP reduces to an LP in standard form.

- Let $g(x) = I_+(x)$ be the indicator function of the positive orthant \mathbf{R}_+^n , and let $f(x) = (1/2) x^T P x + q^T x$, with $\mathbf{dom} f = \{x \mid A x = b\}$.
- The QP can be expressed in ADMM form as

$$\min_{x,u} f(x) + g(u) \quad \text{s.t.} \quad x - u = 0$$

- Consequently, the resulting ADMM iterations are

$$\begin{aligned} x^{(t+1)} &= \arg \min_{x: Ax=b} \{f(x) + (\rho/2) \|x - u^{(t)} + p^{(t)}\|_2^2\} \\ u^{(t+1)} &= (x^{(t+1)} + p^{(t)})_+ \\ p^{(t+1)} &= p^{(t)} + x^{(t+1)} - u^{(t+1)} \end{aligned}$$

- The x -update is an equality-constrained LS problem. Its optimality conditions can be expressed as

$$\overbrace{\begin{bmatrix} P + \rho I & A^T \\ A & 0 \end{bmatrix}}^R \begin{bmatrix} x^{(t+1)} \\ \nu \end{bmatrix} + \begin{bmatrix} q - \rho(u^{(t)} - p^{(t)}) \\ -b \end{bmatrix} = 0$$

where ν is the dual variable related to the equality constraint $Ax = b$.

- To speed up computations, the matrix R can be ("pre") factorized.
- If P is diagonal, possibly zero, this update has a cost of $\mathcal{O}(nm^2)$ flops. Subsequent updates only cost $\mathcal{O}(nm)$ flops.
- The update in u involves projection onto \mathbf{R}_+^n .
- If the non-negativeness constraint $x \succeq 0$ is replaced by a more general *conic* constraint $x \in \mathcal{K}$, where \mathcal{K} is a *convex cone*, the ADMM update for u becomes

$$u^{(t+1)} = \mathcal{P}_{\mathcal{K}}(x^{(t+1)} + p^{(t)})$$

EXAMPLE: IMAGE RECONSTRUCTION

- Let $g \in \mathbf{R}^{m \times n}$ be a *discrete image* which we consider to be a *noisy and blurred* version of some “true image” $f \in \mathbf{R}^{m \times n}$. Formally,

$$g \approx \mathfrak{A}(f)$$

where $\mathfrak{A} : \mathbf{R}^{m \times n} \rightarrow \mathbf{R}^{m \times n}$ is the linear operator representing the effect of blur (i.e., low-pass filter).

- Given g , our task is to find an estimate f . To this end, we consider

$$\min_f \{ \|\mathfrak{A}(f) - g\|_1 + \gamma \|f\|_{\text{TV}} \}$$

with some *regularization parameter* $\gamma > 0$.

- The first term is a *data fidelity* term. It requires the optimal solution to be *close* to g , as measured by the *sum-of-absolute-values* norm $\|\cdot\|_1$.
- The second term is a *prior* term. It requires the solution to have relatively low *variability*, as measured by the *total-variation norm* $\|\cdot\|_{\text{TV}}$.

EXAMPLE: IMAGE RECONSTRUCTION (CONT.)

- Equivalently, the desired estimate can be found as a solution to

$$\begin{aligned} \min_{f,u,v} \quad & \frac{1}{\gamma} \|u - g\|_1 + \|v\|_{\text{TV}} \\ \text{s.t.} \quad & \mathfrak{A}(f) - u = 0, \quad f - v = 0 \end{aligned}$$

- In this case, the ADMM update in x is found as a solution to

$$(\rho_u \mathfrak{A}^* \mathfrak{A} + \rho_v \mathbf{I})(x^{(t+1)}) = \rho_u \mathfrak{A}^*(u^{(t)} - p_u^{(t)}) + \rho_v (v^{(t)} - p_v^{(t)})$$

which is a *linear operator equation* (with \mathfrak{A}^* and \mathbf{I} being the adjoint and identity operators, respectively).

- The ADMM update in u is given by

$$u^{(t+1)} = \mathcal{S}_{1/\rho_u}(\mathfrak{A}(x^{(t+1)}) + p_u^{(t)} + g) - g$$

where \mathcal{S}_{1/ρ_u} is the operator of (coordinate-wise) *soft-thresholding*.

- The ADMM update in v is given by

$$v^{(t+1)} = \mathcal{P}_{\text{TV}^*}(x^{(t+1)} + p_v^{(t)})$$

where $\mathcal{P}_{\text{TV}^*}$ is the operator of projection onto the unit ball in the *dual space* with norm $\|\cdot\|_{\text{TV}^*}$.

EXAMPLE: IMAGE RECONSTRUCTION (CONT.)



original



noisy/blurred



restored

- 1024×1024 image
- Gaussian blur model
- "Salt-and-pepper" noise
- Convergence in less than 500 iterations to $\|\nabla L_{\rho_u, \rho_u}\| \leq 10^{-5}$.

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- 2 web.stanford.edu/~boyd/papers/pdf/admm_distr_stats.pdf
- 3 www.seas.ucla.edu/~vandenbe/publications/pdsplitting.pdf
- 4 A. Chambolle, "An Algorithm for Total Variation Minimization and Applications," *Journal of Mathematical Imaging and Vision*, 20, 89-97, 2004